

Week 3

Students are advised to refer the class discussion for further understanding

Unit 3: Part II: Divide and Conquer

- **Principle:**

- Input is divided into small groups before processing using logical division.
- Solution space for all the divided groups is same and unique.
- The CPU utilization is on higher side whereas the throughput is on lower side.
- The division of input reduces the complexity of the process and most of DAC based algorithms have logarithmic complexity.
- DAC principle is mostly applied for input of large size.

Binary Search

- Principle:

- Implemented on the SORTED array
- Input is divided into small groups before processing using formula:

$$mid = \frac{low+high}{2} \quad (\text{Consider lower integer value})$$

- Where: “low” represents start of array and “high” represents end of array.
- If element is found at position “mid” search is successful, otherwise compare the element at “mid” and element to be searched.
- If $a[mid] < elem \rightarrow low = mid + 1$
- If $a[mid] > elem \rightarrow high = mid - 1$
- Continue process till element is found ($a[mid] = elem$)
- Or otherwise if ($low > high$) \rightarrow Search is unsuccessful.

Binary Search: Numerical on average number of comparisons for successful and unsuccessful search

- Consider the following array of size: 14 (1..14 → low=1 and high=14)

I	1	2	3	4	5	6	7	8	9	10	11	12	13	14
A	-10	-5	7	14	21	34	45	56	65	77	89	98	104	110
C	3	4	2	4	3	4	1	4	3	4	2	4	3	4

- Total number of comparisons for successful search: 43
- Average number of comparisons = $43/14 = 3.07$
- For finding total number of comparisons for unsuccessful search:
 - Draw Binary search tree for the given array
 - Complete the BST and find the sum of level of BST
 - Compute average by dividing the sum/n+1 (where n = number of elements)

Binary Search: Algorithm and complexity equation

- **Algorithm *BinarySearch(a,i,l,x)* //x is the element to be searched**

```
{
    If(i==l) then
    {
        If(x==a[i])
            return(i);
        else
            return(0);
    }

    Else
    {
        Mid=[i+l]/2
        If(x==a[mid]) then
            Return(mid)
        Else
        If(x<a[mid])then
            Return (bsearch(a,i,mid-1,x))
        Else
            Return (bsearch(a,mid+1,l,x));
    }
}
```

Ternary Search

Ternary Search:

Algorithm ternarySearch(a, low, high, searchValue)

```
{  
  
    int left = 1;  
    int right = n;  
    while(left < right)  
    {  
        int middle1 = (right + left) / 3;  
        int middle2 = (right + left) / 3 * 2;  
  
        if(a[middle1] == searchValue)  
        {  
            return middle1;  
        }  
        else if(a[middle2] == searchValue)  
        {  
            return middle2;  
        }  
        else if (searchValue < a[middle1])  
        {  
            right = middle1 - 1;  
        }  
        else if (searchValue > a[middle1] && searchValue < a[middle2])  
        {  
            left = middle1 + 1;  
            right = middle2 - 1;  
        }  
        else if (searchValue > a[middle2])  
            left = middle2;  
    }  
}
```

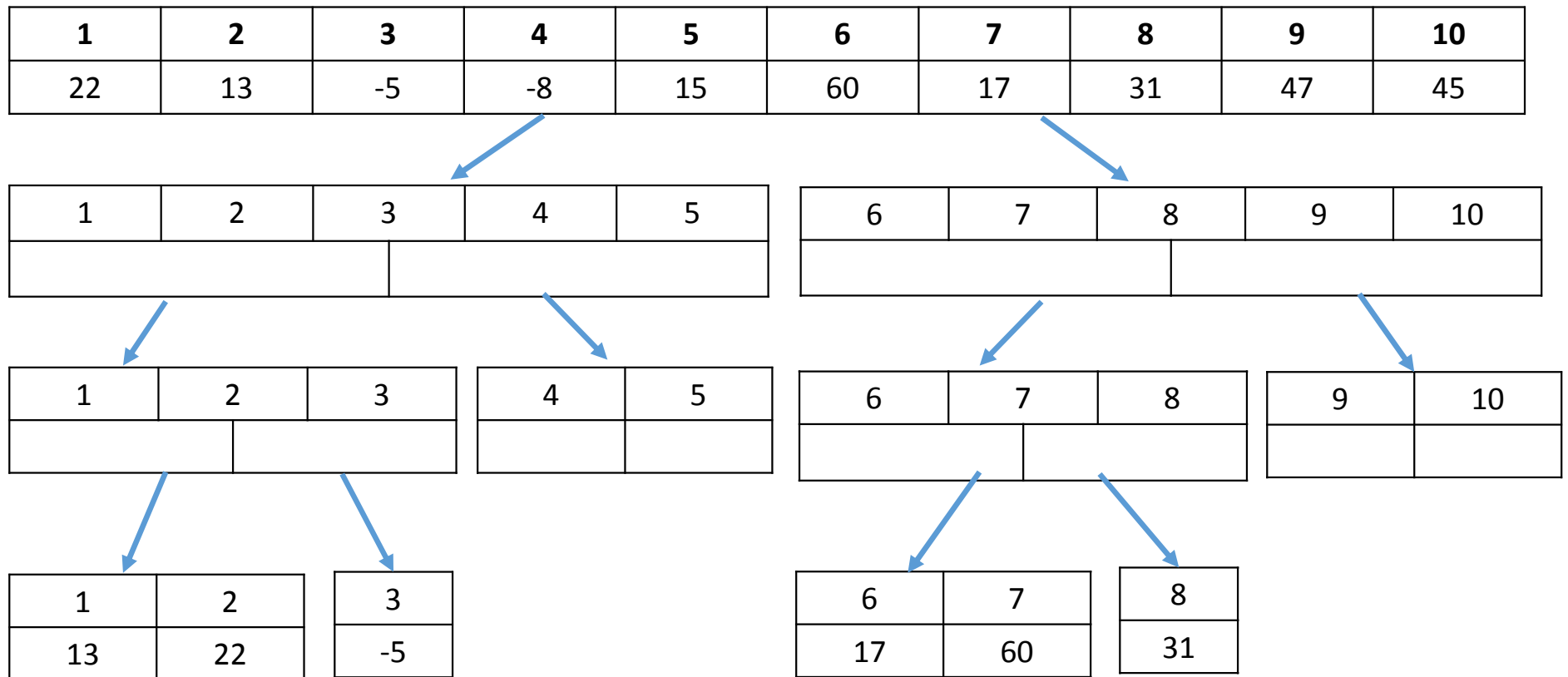
Min-Max Algorithm

- **Principle:** Method to find minimum and maximum element from an array.
- General method requires $2*(n-1)$ comparisons.

Algorithm min_max(a,n,min,max)

```
{
    min = max = a[1]
    for i = 2 to n do
    {
        if (a[i] < min) then min = a[i]
        if (a[i] > max) then max = a[i]
    }
}
```

Min-Max: DAC based Example



Min-Max: DAC Algorithm

Algorithm min_max(a,i,j,min,max)

```
{  
If (i==j) then min=max=a[i]  
Else  
If(i=j-1) then  
{  
  if(a[i] < a[j]) then  
  {  
    min = a[i]; max=a[j]  
  }  
  else  
  {  
    min=a[j]; max=a[i]  
  }  
Else
```

```
{  
  mid = [i+j] / 2 (Take lower integer)  
  min_max(a,i,mid,min,max)  
  min_max(a,mid+1, j, min1, max1)  
  if (max < max1) then max = max1  
  if (min > min1) then min = min 1  
}  
}
```

- Complexity equation:
- $T(n) = T(n/2) + T(n/2) + 2$ for $n > 2$
- $T(n) = 1$ if $n=2$
- $T(n) = 0$ if $n=1$